

# Members of thin $\Pi_1^0$ classes and their Turing degrees

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## $\Pi_1^0$ classes

### Definition:

A  $\Pi_1^0$  class is a set  $P \subseteq 2^\omega$  for which there is a (primitive) recursive tree  $T$  with  $[T] = P$ .

- ▶ As primitive recursive trees can be effectively enumerated, we have an effective enumeration of  $\Pi_1^0$  classes.

### Examples:

- ▶ Consider  $A = \{e : \varphi_e(0) \downarrow = 0\}$  and  $B = \{e : \varphi_e(0) \downarrow = 1\}$ .
  - ▶  $A$  and  $B$  are disjoint r.e. sets, and cannot be recursively separated.
  - ▶ The class of sets separating  $A$  and  $B$

$$S(A, B) = \{C : A \subseteq C \text{ \& } B \cap C = \emptyset\}$$

is called the separating class of  $A$  and  $B$ .

- ▶  $S(A, B)$  is a  $\Pi_1^0$  class and is perfect (hence uncountable).
- ▶ The class of complete consistent extension of Peano Arithmetic is a  $\Pi_1^0$  class.
- ▶ Zariski topology over recursive rings, where for any r.e. ideal  $I$ , the collection of prime ideals containing  $I$  forms a  $\Pi_1^0$  class.

## Thin $\Pi_1^0$ classes

### Definition: Thin Classes

A  $\Pi_1^0$  class  $P$  is thin if every subclass of  $P$  is relatively clopen, i.e., if  $Q$  is a subclass of  $P$ , then  $Q = P \cap U$  for some clopen set  $U \subseteq 2^\omega$ .

We know that in all  $\Pi_1^0$  classes, isolated paths are computable.

Conversely, if a thin  $\Pi_1^0$  class  $P$  contains a computable element  $X$ , then  $\{X\}$  is a subclass of  $P$ , and hence by the thinness of  $P$ ,  $X$  is isolated.

### FACT:

A thin  $\Pi_1^0$  class  $P$  has no computable members if and only if  $P$  is perfect.

So every **countable thin  $\Pi_1^0$  class** has a computable member.

## Martin-Pour El theories

The notion of thinness comes from the work of Martin and Pour-El in 1970. Let  $S$  be a consistent r.e. theory in the propositional language with

### Martin and Pour-El, 1970

Let  $S$  be a consistent r.e. theory.

- (1)  $S$  has few r.e. extensions if each r.e. extension  $T$  of  $S$  is a principal extension, i.e.,  $T$  is generated by  $S$  together with a single propositional formula.
- (2)  $S$  is essentially undecidable if  $S$  has no decidable complete consistent extensions.

### FACTS:

For a consistent r.e. theory  $S$ ,

- ▶  $S$  has few r.e. extensions if and only if the corresponding  $\Pi_1^0$  class is thin.
- ▶  $S$  is essentially undecidable if and only if the corresponding  $\Pi_1^0$  class has no computable members.

## Turing degrees of members of thin $\Pi_1^0$ classes

Theorem (CDJS, 1993):

If  $X$  is in a thin  $\Pi_1^0$  class  $P$ , then  $X' \leq_T X \oplus \emptyset''$ .

**Proof:** Let  $P = [T]$  is a thin class, where  $T$  is a recursive tree, and  $A \in P$ .

For a given  $e$ , we consider whether  $e \in A'$  or not, i.e., whether  $\Phi_e^A(e) \downarrow$  or not.

- ▶ If  $\Phi_e^A(e) \downarrow$ , we can recursive in  $A$  to find an initial segment  $\sigma$  of  $A$  with  $\{e\}^\sigma(e) \downarrow$ .
- ▶ If NOT, what shall we do?

Consider  $Q_e = \{C : \Phi_e^C(e) \uparrow\}$ , a  $\Pi_1^0$  class

- ▶  $P \cap Q_e$  is a subclass of  $P$ , and as  $P$  is thin,  $P \cap Q_e = P \cap U_e$  for some clopen set  $U_e$ .
- ▶ As we are assuming that  $A$  is in  $P \cap Q_e$ ,  $A$  is in  $P \cap U_e$ , and hence  $A$  has an initial segment  $\sigma$  with all infinite extensions in  $U_e$ .

Thus, if  $B \in P$  extends  $\sigma$ , then  $B \in P \cap U_e = P \cap Q_e$ , and  $\Phi_e^B(e) \uparrow$ .

- ▶ Define a binary relation  $R(e, \sigma)$  as

$$R(e, \sigma) \iff (\forall \tau \supseteq \sigma)[\tau \in T \ \& \ \{e\}^\tau(e) \downarrow \rightarrow \tau \notin \text{Ext}(T)].$$

$R$  is a  $\Pi_2$  relation and is recursive in  $\emptyset''$ .

We do as following:

Find the least number  $n$  such the following is true for  $\sigma = A \upharpoonright n$ :

- (a)  $\{e\}^\sigma(e) \downarrow \rightarrow e \in A'$
- (b)  $R(e, \sigma) \rightarrow e \notin A'$

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If  $A$  computes  $\emptyset''$ , then  $A$  cannot be a member of any thin  $\Pi_1^0$  class.

## Spector's construction

- ▶ In Spector's construction of minimal degrees below  $\mathbf{0}''$ , forcing notions are recursive perfect trees,  $T_e$ ,  $e \in \omega$ , pruned according to [the black-white rule](#).

That is, to see whether we can find a string  $\sigma \in T_e$  such that there is no  $e$ -splitting above  $\sigma$  in  $T_e$ , or not.

- ▶ If we use only 'half' of each  $T_e$ , i.e., keep the even part, and exclude the odd part, the construction still works.



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- ▶ If we use only 'half' of each  $T_e$ , i.e., keep the even part, and exclude the odd part, the construction still works.

A great observation.

## Construct one thin-free degree below $\mathbf{0}''$

### Definition:

A Turing degree is thin-free, if no members in this degree is a member of thin  $\Pi_1^0$  classes.

Note that all degrees above  $\mathbf{0}''$  are thin-free.

We will construct a set  $A$  of thin-free degree below  $\mathbf{0}''$ , we shall ensure for any  $e$  such that if  $\Phi_e^A$  is total and Turing equivalent to  $A$ , then one of the following is guaranteed:

- (1)  $\Phi_e^A \notin [P_e]$ , or
- (2)  $[P_e]$  is not thin.

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Suppose that  $A$  is constructed on a given recursive perfect tree  $T$ .

- ▶ To meet (1), we try to find some string  $\tau$  on  $T$  such that  $\Phi_e^\tau$  is not extendible on  $P_e$ ,
  - ▶ if such a  $\tau$  exists, we force  $A$  to extend  $\tau$ , which guarantees that  $\Phi_e^A \notin [P_e]$ , if  $\Phi_e^A$  is total.

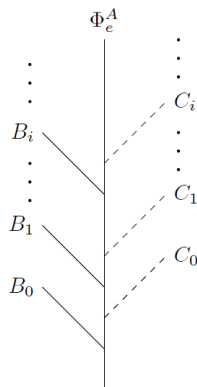
If NOT, we will then try to

- ▶ find a  $\Pi_1^0$  subclass of  $[P_e]$  which is not the intersection of  $[P_e]$  with any clopen set  $U$ .

We will construct a recursive subtree  $S_e$  of  $P_e$ , such that  $\Phi_e^A$  lies on  $S_e$ , and for any length  $n$ , there exists some  $B \in [S_e]$  and  $C \in [P_e] \setminus [S_e]$  such that

$$B \upharpoonright n = C \upharpoonright n = \Phi_e^A \upharpoonright n.$$

This implies that  $\Phi_e^A \in [P_e]$ , and  $[S_e]$  witnesses that  $P_e$  is not thin.



## Action under this case:

- ▶ **Target:** Force  $A$  on a total recursive subtree  $T_e$  of  $T$ , such that for any  $\alpha \in T_e$ ,  $\Phi_e^{T_e(\alpha 0)}$  and  $\Phi_e^{T_e(\alpha 1)}$  are incompatible in  $P_e$  and there is a path on  $P_e$  extending  $\Phi_e^{T_e(\alpha)}$ , of course.

*We are assuming that (1) fails, so both  $\Phi_e^{T_e(\alpha 0)}$ ,  $\Phi_e^{T_e(\alpha 1)}$  are extendible on  $P_e$  and thus there is at least one infinite path in  $P_e$  extending it.*

Consider the  $e$ -splitting subtree of  $T$ ,  $SP(T, e)$ , if exists, and take the even part.

- ▶ White Side:  $SP(T, e)$  exists.

In this case,  $\Phi_e^A$  is total, then  $\Phi_e^A$  is on  $[P_e]$ , and  $E(SP(T, e))$ , the even subtree of  $SP(T, e)$ , is a total recursive subtree of  $T$ , and  $\Phi_e^{E(SP(T, e))}$  is a total recursive subtree of  $P_e$ , witnessing that  $[P_e]$  is not thin.

- ▶ Black Side:  $SP(T, e)$  does not exist.

In this case, there is a string  $T(\alpha)$  such that above  $T(\alpha)$ , no string  $e$ -splits, and hence, if  $A$  is on the full subtree of  $T$  above  $\alpha$ ,  $Full(T, \alpha)$ , then  $\Phi_e^A$  is recursive, making  $A$  and  $\Phi_e^A$  not Turing equivalent, if we can make  $A$  nonrecursive. **We Can**, as recursive sets are all in thin  $\Pi_1$  classes..

## Oracle Construction:

We can now run a forcing argument to construct  $A$  with wanted property.

- ▶  $\mathbf{0}''$  is used as oracle to make decision at every stage.

Yuan Bowen improved this in his thesis:

### Theorem:

There exists a hyperimmune-free minimal degree below  $\mathbf{0}''$  which is also thin-free.

Note that such degrees are not below  $\mathbf{0}'$ .

## Working below $\mathbf{0}'$

- ▶ CDJS proved that  $\mathbf{0}'$  contains a  $\Pi_1$  set  $A$  which is in a thin  $\Pi_1$  class  $P$ .
- ▶ CDJS proved the density of degrees containing sets (not necessarily r.e.) in thin  $\Pi_1$  classes in r.e. degrees.

DWY strengthened this in 2018, showing that sets above can be r.e.

- ▶ Yuan Bowen proved in his thesis that all 1-generic degrees below  $\mathbf{0}'$  contain members of thin  $\Pi_1$  classes.



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- ▶ Yuan Bowen proved in his thesis that all 1-generic degrees below  $\mathbf{0}'$  contain members of thin  $\Pi_1$  classes.
- ▶ There are degrees below  $\mathbf{0}'$  thin-free and then can be r.e., or minimal, by CDJS.

The construction of a minimal thin-free degree was given by CDJS, modified from Sacks forcing, where partial recursive trees are used.

## An r.e. thin-free degree

Construct an r.e.  $A$  satisfying the following requirements:

$\mathcal{R}_e$ : if  $\Phi_e(A)$  and  $\Psi_e(\Phi_e(A))$  are both total, then either

- ▶  $A \neq \Psi_e(\Phi_e(A))$ ; or
- ▶  $\Phi_e(A)$  is not in  $[P_e]$ ; or
- ▶  $[P_e]$  is not thin.

In this construction, we cannot use the  $e$ -splitting tree as a help to construct a subclass witness that  $[P_e]$  is not thin.

We thus need to construct such a subclass, actually, a subtree, by infinitely many substrategies, each of which tries to find an infinite path in  $[P_e]$ , and

- ▶ any substrategy fails to secure an infinite path, an enumeration of a certain number into  $A$ , showing that either  $A \neq \Psi_e(\Phi_e(A))$  (diagonalization succeeds) or  $\Phi_e(A)$  is not in  $[P_e]$ , a global win for  $\mathcal{R}_e$ .

DWY proved in 2018 that such r.e. degrees are dense in the r.e. degrees.

## Other topics

In his thesis, Yuan proved that any nonrecursive set below a 2-generic set is thin-free. In particular, 2-generic degrees are thin-free.

CDJS also consider minimal  $\Pi_1^0$  classes and Cantor-Bendixson rank of sets, a topic originated from Cenzer, et al.'s work in 1986.

Our continuing work on this topic is in the direction of Ershov hierarchy, also 1-generic degrees not below  $0'$ ,  $pb$ -generic degrees, minimal degrees with full approximations.

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3. Downey, Wu and Yang, The members of thin and minimal  $\Pi_1^0$  classes, their ranks and Turing degrees, Annals of Pure and Applied Logic **166** (2015), 755-766.
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5. Yuan Bowen, PhD thesis, NTU, 2020.

**Thanks!**

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**Take care and keep safe!**