

# HANDOUT

## Singapore Logic Seminar, 03-21-2021

Speaker: Lars Kristiansen

This talk is based on material published in [1, 2, 3, 4] and some unpublished material.

*Do we need, or do we not need, unbounded search in order to convert one representation of an irrational number into another representation?*

⋮

*Do we need, or do we not need, unbounded search in order to convert a Cauchy sequence for  $\alpha$  into the Dedekind cut of  $\alpha$ ?*

*Do we need, or do we not need, unbounded search in order to convert the Dedekind cut of  $\alpha$  into a Cauchy sequence for  $\alpha$ ?*

*... the base-2 expansion of  $\alpha$  into the base-10 expansion of  $\alpha$ ?*

*... the base-10 expansion of  $\alpha$  into the base-2 expansion of  $\alpha$ ?*

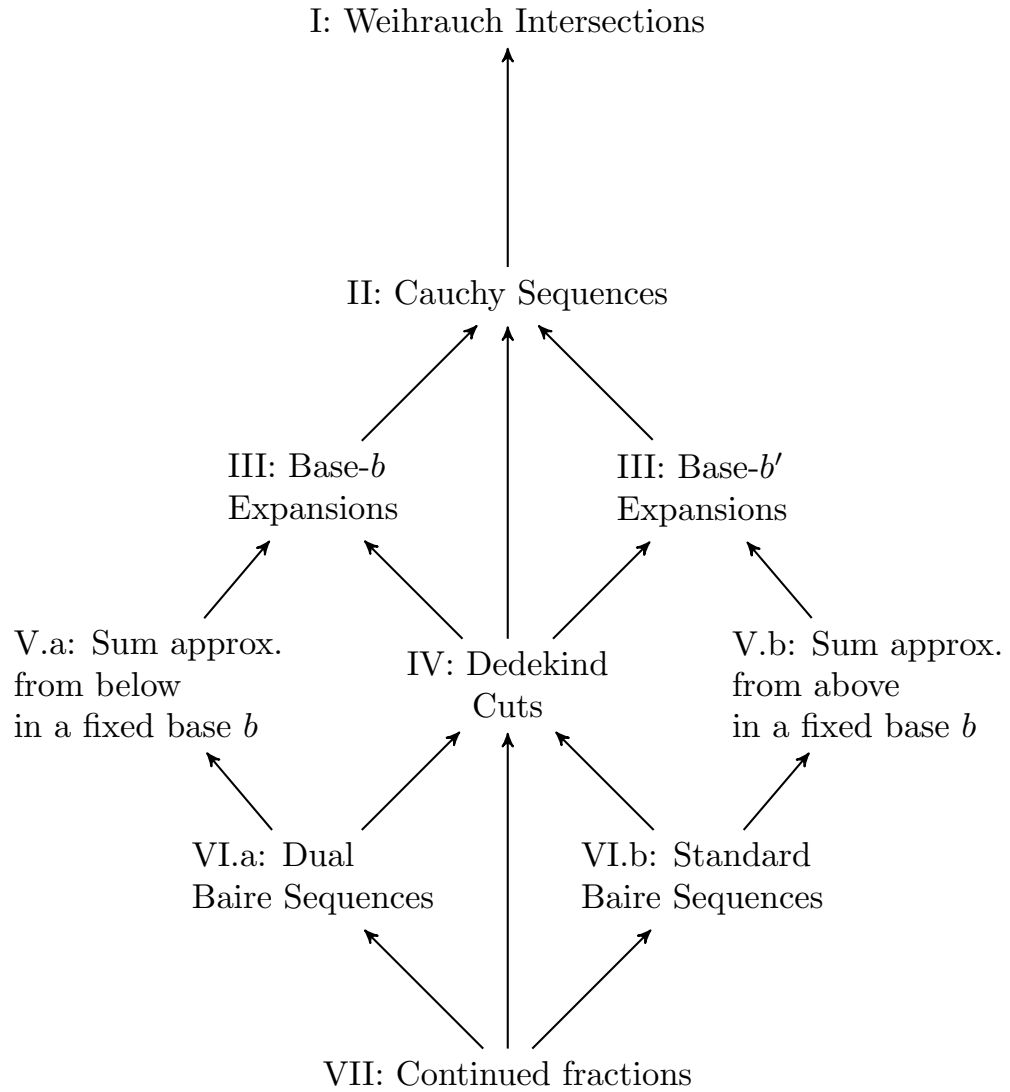
*... the continued fraction of  $\alpha$  into the base-17 expansion of  $\alpha$ ?*

*... the Dedekind cut of  $\alpha$  into the continued fraction of  $\alpha$ ?*

⋮

A computation that does not apply unbounded search is called a *subrecursive* computation. Primitive recursive computations and (Kalmar) elementary computations are typical examples of subrecursive computations. A representation  $R_1$  (of irrational numbers) is *subrecursive in* a representation  $R_2$  if the  $R_1$ -representation of  $\alpha$  can be subrecursively computed in the  $R_2$ -representation of  $\alpha$ . Two representations  $R_1$  and  $R_2$  are *equivalent* when  $R_1$  is subrecursive in  $R_2$  and  $R_2$  is subrecursive in  $R_1$ .

Overview of subrecursive reducibility among  
(equivalence classes of) representations:



## I: Representations Equivalent to Weihrauch Intervals

Let  $\alpha$  be an irrational number in the interval  $(0, 1)$ .

### Weihrauch Intervals

A function  $I : \mathbb{N} \rightarrow \mathbb{Q} \times \mathbb{Q}$  is a *Weihrauch intersection* of  $\alpha$  if

$$\{ \alpha \} = \bigcap_n I_n^O$$

where  $I_n^O$  denotes the open interval given by the pair  $I(n)$ .

*Comments:* This is a representation from Weihrauch's book [19].

### Nested Weihrauch Intervals

A *nested Weihrauch intersection*  $I : \mathbb{N} \rightarrow \mathbb{Q} \times \mathbb{Q}$  is a Weihrauch intersection such that  $I_{n+1}^O$  is a strict subinterval of  $I_n^O$ .

### Complete Topological Names

Let  $f : \mathbb{N} \rightarrow \mathbb{Q} \times \mathbb{Q}$  be such that (1) for any open interval  $I$  with rational endpoints and  $\alpha \in I$  exists  $n$  such that  $f(n) = (r_1, r_2) = I$  and (2) if  $x \notin I$ , we have  $f(n) \cap I = \emptyset$  for all  $n$ . So  $\{f(i)\}_{i \in \mathbb{N}}$  is a sequence whose elements are exactly the open intervals with rational endpoints that contains  $\alpha$ , and we have  $\{\alpha\} = \bigcap_i f(i)$ . Then we say that  $f$  is a *complete topological name* for  $\alpha$ .

*Comments:* This is a representation from Weihrauch's book [19].

## II: Representations Equivalent to Cauchy Sequences

Let  $\alpha$  be an irrational number in the interval  $(0, 1)$ .

### Cauchy Sequences

The function  $C : \mathbb{N} \rightarrow \mathbb{Q}$  is a *Cauchy sequence* for  $\alpha$  if

$$|\alpha - C(n)| < \frac{1}{2^n}.$$

### Strictly Increasing Cauchy Sequences

The function  $C : \mathbb{N} \rightarrow \mathbb{Q}$  is a *strictly increasing Cauchy sequence* for  $\alpha$  if (i)  $C$  is a Cauchy sequence for  $\alpha$  and (ii)  $C(n) < C(n+1)$ .

### Base- $b$ Cauchy Sequences

Let  $b \in \mathbb{N} \setminus \{0, 1\}$ . The function  $f : \mathbb{N} \rightarrow \mathbb{N}$  is a *base- $b$  Cauchy sequence* for  $\alpha$  if

$$C(n) := \frac{f(n)}{b^n}$$

is a Cauchy sequence for  $\alpha$ .

*Comments:* Friedman & Ko [8] use base-2 Cauchy sequences. They call them “converging sequences of dyadic rational numbers”. *Increasing base- $b$  Cauchy sequences* will not be equivalent to base- $b$  Cauchy sequences.

### Fuzzy (Dedekind) Cuts

The function  $D : \mathbb{N} \times (\mathbb{N} \setminus \{0\}) \rightarrow \{0, 1\}$  is a *fuzzy Dedekind cut* for  $\alpha$  if

$$D(m, n) = 0 \Rightarrow \alpha < \frac{m+1}{n} \quad \text{and} \quad D(m, n) = 1 \Rightarrow \frac{m-1}{n} < \alpha.$$

### Signed Digit Expansions

The function  $S : (\mathbb{N} \setminus \{0, 1\}) \rightarrow \{-1, 0, 1\}$  is a *signed digit expansion* of  $\alpha$  if  $S(0) = 0$  and

$$\alpha = \sum_{i=1}^{\infty} \frac{S(i)}{2^i}.$$

*Comments:* This seems to be a pretty standard representation, see Berger et al. [7].

### III: Equivalence Classes of Base- $b$ Expansions

Let  $\alpha$  be an irrational number in the interval  $(0, 1)$ .

The function  $f : \mathbb{N} \rightarrow \{0, \dots, b-1\}$  is the *base- $b$  expansion* of  $\alpha$  if  $E(0) = 0$  and

$$\alpha = \sum_{i=1}^{\infty} \frac{f(i)}{b^i}.$$

We use  $E_b^\alpha$  to denote the base- $b$  expansion of  $\alpha$ .

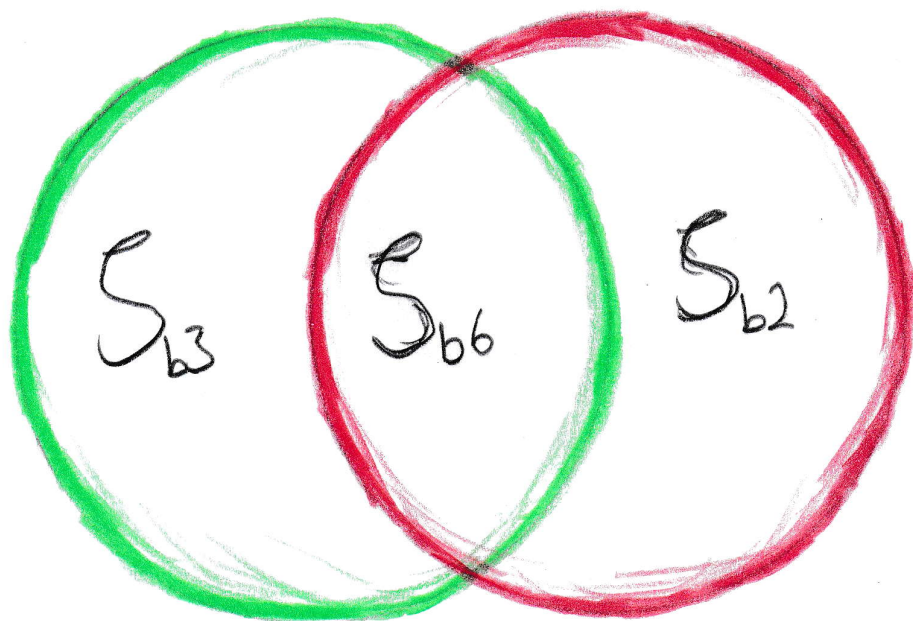
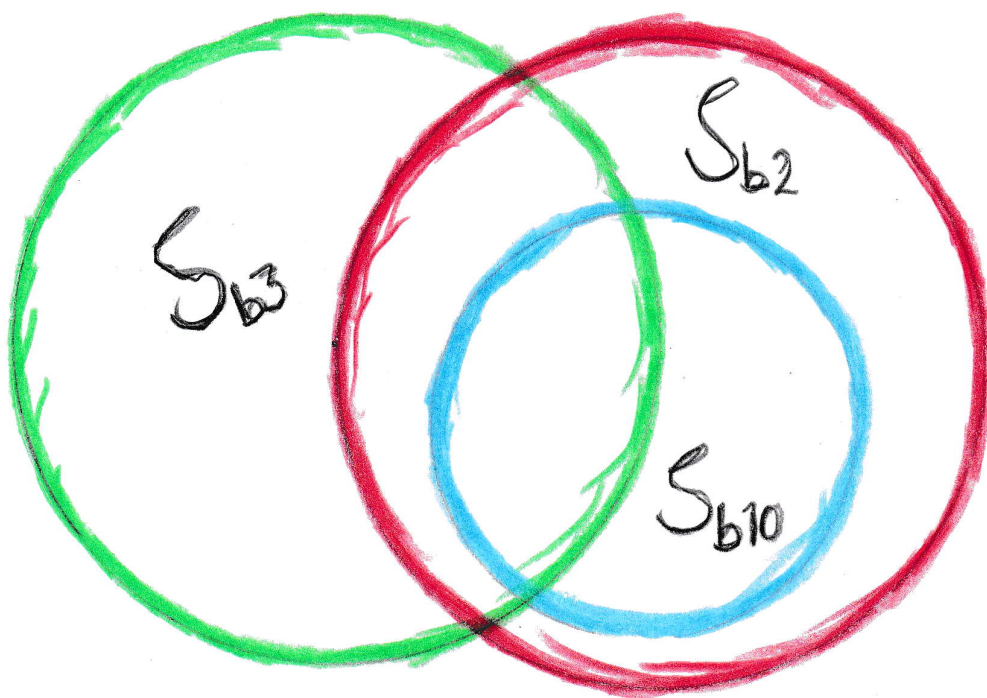
The next theorem implies that there is a lot of degrees between the degree of the Cauchy sequences and the degree of the Dedekind cuts.

**Theorem 1** *The base- $b$  expansion is subrecursive in the base- $b'$  expansion if and only if every prime that divides  $b$  also divides  $b'$ .*

*Comments:* Every prime that divides  $b$  will also divide  $b'$  if and only if every rational number that has a finite base- $b$  expansion also has a finite base- $b'$  expansion.

We may also consider Venn diagrams. In the Venn diagrams AT THE NEXT PAGE we use the notation:

- $\mathcal{S}$  may be any subrecursive class closed under elementary operation
- $\mathcal{S}_{bn}$  is the class of irrational numbers that have a base- $n$  expansion in  $\mathcal{S}$ .



## IV: Representations Equivalent to Dedekind Cuts

Let  $\alpha$  be an irrational number in the interval  $(0, 1)$ .

### Dedekind Cuts

The function  $D : \mathbb{Q} \rightarrow \{0, 1\}$  is the *Dedekind cut* of  $\alpha$  if

$$D(q) = \begin{cases} 0 & \text{if } q < \alpha \\ 1 & \text{if } \alpha < q. \end{cases}$$

### General Base Expansions

The function  $E : (\mathbb{N} \setminus \{0, 1\}) \times \mathbb{N} \rightarrow \{0, \dots, b-1\}$  is the *general base expansion* of  $\alpha$  if  $E(b, n) = E_b^\alpha(n)$  (the base- $b$  expansion  $E_b^\alpha$  is defined above above).

### Beatty Sequences

The function  $B : (\mathbb{N} \setminus \{0\}) \rightarrow \mathbb{N}$  is the *Beatty sequence* of  $\alpha$  if

$$\frac{B(n)}{n} < \alpha < \frac{B(n) + 1}{n}.$$

*Comments:* The name *Beatty sequence* has its origin in the publication [5]. Apparently, what is now known as Beatty sequences was used earlier by Bernard Bolzano [6], whence this representation could also be called *Bolzano measures*.

### Hurwitz Characteristics

For any string  $\tau \in \{L, R\}^*$ , we define the *interval addressed by  $\tau$*  inductively over the structure of  $\tau$ : The empty sequence addresses the interval  $(0/1, 1/1)$ . Furthermore

$$\tau L \text{ addresses } \left( \frac{a}{b}, \frac{a+c}{b+d} \right) \text{ and } \tau R \text{ addresses } \left( \frac{a+c}{b+d}, \frac{c}{d} \right)$$

if  $\tau$  addresses  $(a/b, c/d)$ .

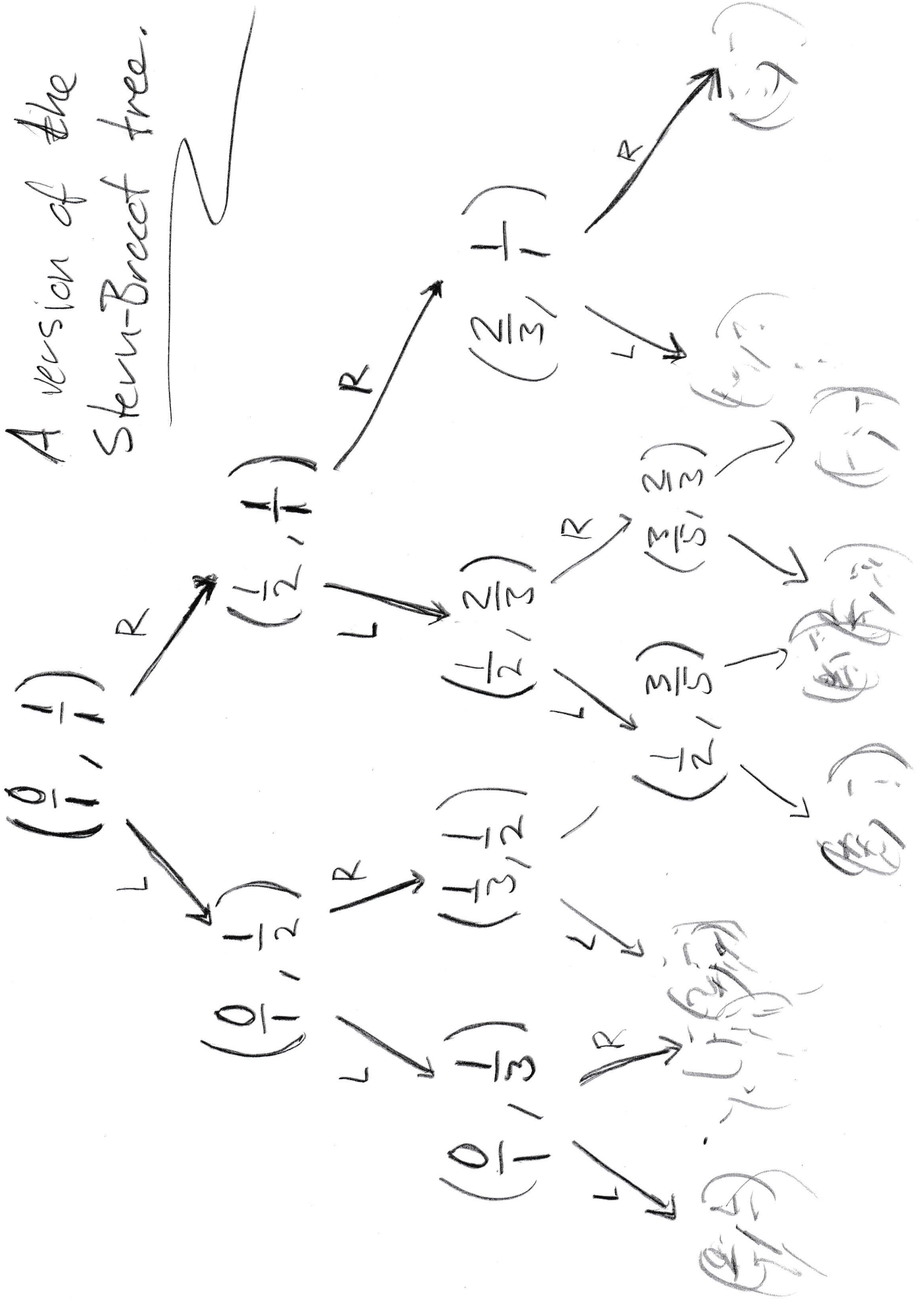
The infinite sequence  $\Sigma$  over the alphabet  $\{L, R\}$  addresses  $\alpha$  if any finite prefix of  $\Sigma$  addresses an interval that contains  $\alpha$ .

Let  $\Sigma^\alpha$  address  $\alpha$ . The function  $H : (\mathbb{N} \setminus \{0\}) \rightarrow \{0, 1\}$  is the *Hurwitz characteristic* of  $\alpha$  if

$$H(n) = \begin{cases} 0 & \text{if the } n\text{'th element of } \Sigma^\alpha \text{ is } L \\ 1 & \text{if the } n\text{'th element of } \Sigma^\alpha \text{ is } R. \end{cases}$$

*Comments:* Hurwitz characteristics were known in the 19th century, see Hurwitz [10]. For more on Hurwitz characteristics (as representation of irrationals) see Lehman [14] and Kristiansen & Simonsen [4]. A Hurwitz characteristic yields a branch in the Stern-Brocot tree, see the figure AT THE NEXT PAGE.

A version of the Stern-Brocot tree.





## V: Equivalence Classes of Sum Approximations

Let  $\alpha$  be an irrational number in the interval  $(0, 1)$ .

Let  $0.D_1D_2\dots$  be the base- $b$  expansion of  $\alpha$  (and recall that  $E_b^\alpha(i) = D_i$ ). If  $D$  is a base- $b$  digit, then  $\bar{D}$  denotes the *complement digit* of  $D$ , that is,  $\bar{D} = (b-1) - D$ .

The *base- $b$  sum approximation from below* of  $\alpha$  is the function  $\hat{A}_b^\alpha : \mathbb{N} \rightarrow \mathbb{Q}$  defined by  $\hat{A}_b^\alpha(0) = 0$  and

$$\hat{A}_b^\alpha(n+1) = \frac{E_b^\alpha(m)}{b^m}$$

where  $m$  is the least  $m$  such that

$$\sum_{i=0}^n \hat{A}_b^\alpha(i) < 0.D_1\dots D_m.$$

The *base- $b$  sum approximation from above* of  $\alpha$  is the function  $\check{A}_b^\alpha : \mathbb{N} \rightarrow \mathbb{Q}$  defined by  $\check{A}_b^\alpha(0) = 0$  and

$$\check{A}_b^\alpha(n+1) = \frac{\overline{E_b^\alpha(m)}}{b^m}$$

where  $m$  is the least  $m$  such that

$$1 - \sum_{i=0}^n \check{A}_b^\alpha(i) > 1 - 0.\bar{D}_1\dots\bar{D}_m.$$

Observe that we have

$$\sum_{i=0}^{\infty} E_b^\alpha(i) = \sum_{n=0}^{\infty} \hat{A}_b^\alpha(n) = 1 - \sum_{n=0}^{\infty} \check{A}_b^\alpha(n).$$

**Example:** Let the base-10 expansion of  $\alpha$  start with the digits  $0.3000604\dots$ . Then we have

$$\hat{A}_b^\alpha(1) = 3 \times 10^{-1} \quad \hat{A}_b^\alpha(2) = 6 \times 10^{-5} \quad \hat{A}_b^\alpha(3) = 4 \times 10^{-7}$$

and

$$\check{A}_b^\alpha(1) = 6 \times 10^{-1} \quad \check{A}_b^\alpha(2) = 9 \times 10^{-2} \quad \check{A}_b^\alpha(3) = 9 \times 10^{-3}.$$

**Exercise:** Assume that the base- $b$  expansion of  $\alpha$  contains very long sequences of zeros. Describe the base- $b$  sum approximation from below of  $\alpha$ . Describe the base- $b$  sum approximation from above of  $\alpha$ .

We say that a representation  $R_1$  is *incomparable* to a representation  $R_2$  if (i)  $R_1$  is not subrecursive in  $R_2$  and (ii)  $R_2$  is not subrecursive in  $R_1$ . The next theorems show that there is a lot of degrees that are incomparable to the degree of the Dedekind cut.

**Theorem 2** *Let  $b$  be an arbitrary base (so  $b \geq 2$ ). Then*

- *the base- $b$  sum approximation from below is incomparable to the base- $b$  sum approximation from above*
- *the base- $b$  sum approximation from below is incomparable to the Dedekind cut*
- *the base- $b$  sum approximation from above is incomparable to the Dedekind cut.*

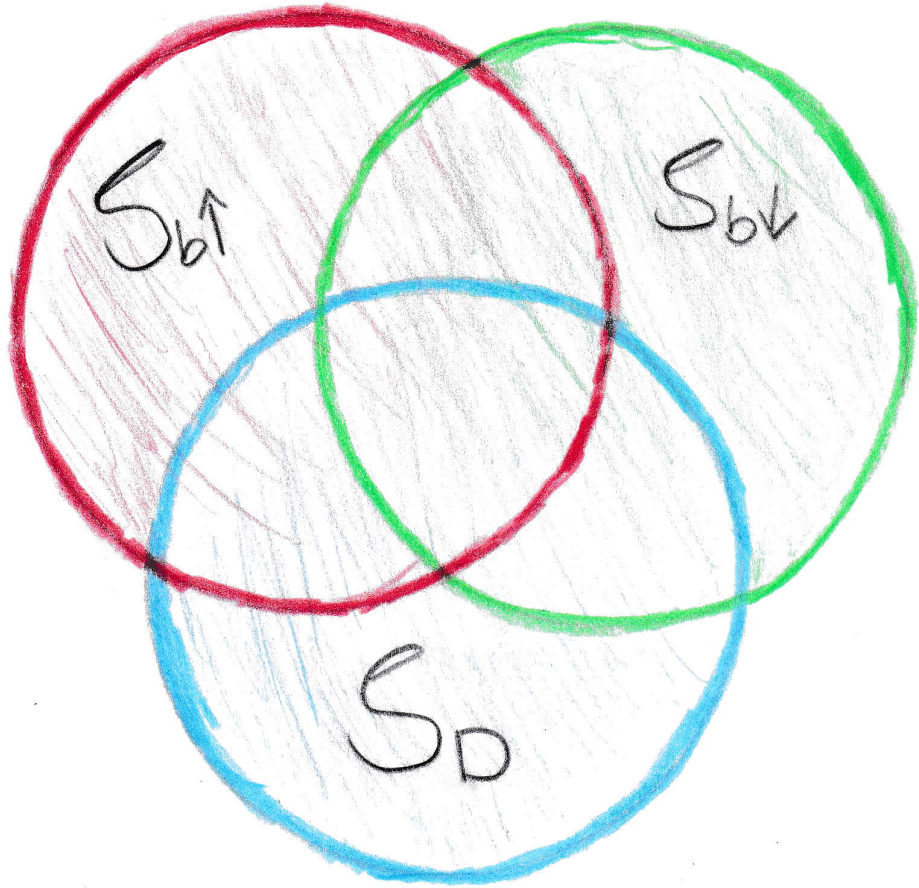
**Theorem 3** *The base- $b$  sum approximation from below of  $\alpha$  is subrecursive in the base- $b'$  sum approximation from below of  $\alpha$  if and only if every prime that divides  $b$  also divides  $b'$ .*

**Theorem 4** *The base- $b$  sum approximation from above of  $\alpha$  is subrecursive in the base- $b'$  sum approximation from above of  $\alpha$  if and only if every prime that divides  $b$  also divides  $b'$ .*

*Comments:* Observe the similarities between the two preceding theorems and Theorem 1.

The Venn diagram AT THE NEXT PAGE gives a little bit more informative than Theorem 2. In the diagram we use the notation:

- $\mathcal{S}$  may be any subrecursive class close under elementary operations
- $\mathcal{S}_D$  is the class of irrational numbers that have a Dedekind cut in  $\mathcal{S}$
- $\mathcal{S}_{b\uparrow}$  is the class of irrational numbers that have base- $b$  sum approximation from below in  $\mathcal{S}$
- $\mathcal{S}_{b\downarrow}$  is the class of irrational numbers that have base- $b$  sum approximation from above in  $\mathcal{S}$



## VI: Representations Equivalent to (Standard and Dual) Baire Sequences

Let  $\alpha$  be an irrational number in the interval  $(0, 1)$ .

### Standard Baire Sequences

Let  $f : \mathbb{N} \rightarrow \mathbb{N}$  be any function, and let  $n \in \mathbb{N}$ . We define the interval  $I_f^n$  by  $I_f^0 = (0/1, 1/1)$  and

$$I_f^{n+1} = \left( \frac{a + f(n)c}{b + f(n)d}, \frac{a + f(n)c + c}{b + f(n)d + d} \right)$$

if  $I_f^n = (a/b, c/d)$ . The function  $B : \mathbb{N} \rightarrow \mathbb{N}$  is the *standard Baire representation* of  $\alpha$  if we have  $\alpha \in I_B^n$  for every  $n$ .

### Dual Baire Sequences

We define the interval  $J_f^n$  by  $J_f^0 = (0/1, 1/1)$  and

$$J_f^{n+1} = \left( \frac{a + f(n)a + c}{b + f(n)b + d}, \frac{f(n)a + c}{f(n)b + d} \right)$$

if  $J_f^n = (a/b, c/d)$ . The function  $A : \mathbb{N} \rightarrow \mathbb{N}$  is the *dual Baire representation* of  $\alpha$  if we have  $\alpha \in J_A^n$  for every  $n$ .

### General Sum Approximations from Below

The *general sum approximation from below* of  $\alpha$  is the function

$$\hat{G}^\alpha : ((\mathbb{N} \setminus \{0, 1\}) \times \mathbb{N}) \rightarrow \mathbb{Q}$$

given by

$$\hat{G}^\alpha(b, x) = \hat{A}_b^\alpha(x)$$

where  $\hat{A}_b^\alpha$  is the base- $b$  sum approximation from below of  $\alpha$  (see definition above).

### General Sum Approximations from Above

The *general sum approximation from above* of  $\alpha$  is the function

$$\check{G}^\alpha : ((\mathbb{N} \setminus \{0, 1\}) \times \mathbb{N}) \rightarrow \mathbb{Q}$$

given by

$$\check{G}^\alpha(b, x) = \check{A}_b^\alpha(x)$$

where  $\check{A}_b^\alpha$  is the base- $b$  sum approximation from above of  $\alpha$  (see definition above).

### Left Best Approximations

Let  $a$  and  $b$  be relatively prime natural numbers with  $b > 0$ . The fraction  $a/b$  is a *left best approximant* of  $\alpha$  if we have  $c/d \leq a/b < \alpha$  or  $\alpha < c/d$  for any natural numbers  $c, d$  with  $0 < d \leq b$ . A *left best approximation* of  $\alpha$  is a sequence of fractions  $\{a_i/b_i\}_{i \in \mathbb{N}}$  such that

$$(0/1) = (a_0/b_0) < (a_1/b_1) < (a_2/b_2) < \dots$$

and each  $a_i/b_i$  is a left best approximant of  $\alpha$ .

### Right Best Approximations

Let  $a$  and  $b$  be relatively prime natural numbers with  $b > 0$ . The fraction  $a/b$  is a *right best approximant* of  $\alpha$  if we have  $\alpha < a/b \leq c/d$  or  $c/d < \alpha$  for any natural numbers  $c, d$  with  $0 < d \leq b$ . A *right best approximation* of  $\alpha$  is a sequence of fractions  $\{a_i/b_i\}_{i \in \mathbb{N}}$  such that

$$(1/1) = (a_0/b_0) > (a_1/b_1) > (a_2/b_2) > \dots$$

and each  $a_i/b_i$  is a right best approximant of  $\alpha$ .

**Theorem 5** *The representations*

- *Dual Baire Sequences*
- *Left Best Approximations*
- *General Sum Approximations from Below*

*are equivalent. The representations*

- *Standard Baire Sequences*
- *Right Best Approximations*
- *General Sum Approximations from Above*

*are equivalent. Moreover, any representation in the first equivalence class is incomparable to any representation in the second equivalence class.*

## VII: Representations Equivalent to Continued Fractions

Let  $\alpha$  be an irrational number in the interval  $(0, 1)$ .

*Comments:* Continued fractions are well known from the literature. The continued fraction  $[0; a_1, a_2, \dots]$  of  $\alpha$  is the unique sequence of positive integers such that

$$\alpha = 0 + \frac{1}{a_1 + \frac{1}{a_2 + \frac{1}{a_2 + \dots}}}$$

### Continued Fractions

The function  $f : (\mathbb{N} \setminus \{0\}) \rightarrow (\mathbb{N} \setminus \{0\})$  is the *continued fraction* of  $\alpha \in (0, 1)$  if  $\alpha = [0; f(1), f(2), f(2), \dots]$ .

### Trace Functions

A function  $T : [0, 1] \cap \mathbb{Q} \rightarrow (0, 1) \cap \mathbb{Q}$  is a *trace function* for the  $\alpha$  if

$$|\alpha - q| > |\alpha - T(q)| .$$

### Contractors

A function  $F : [0, 1] \cap \mathbb{Q} \rightarrow (0, 1) \cap \mathbb{Q}$  is a *contractor* if we have

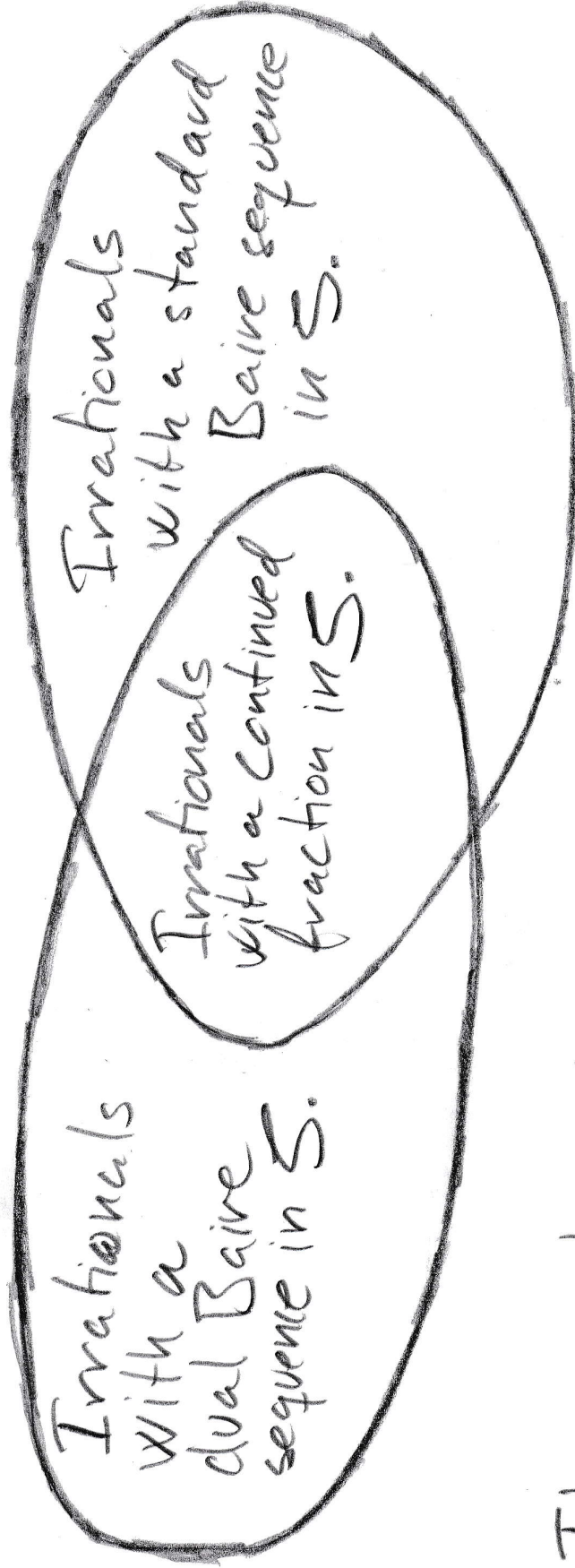
$$F(q) \neq q \quad \text{and} \quad |F(q_1) - F(q_2)| < |q_1 - q_2| .$$

Moreover,  $F$  is a *contractor for*  $\alpha$  if  $F$  is a trace function for  $\alpha$ .

*Comments:* It is easy to prove that any contractor is a trace function for some irrational number.

The Venn diagram at THE NEXT PAGE shows how continued fractions relate to standard and dual Baire sequences.

Let  $S$  be any subrecursive class closed under primitive recursion.



The irrationals that have a continued fraction in  $S$  are exactly those that have both a standard and a dual Baire sequence in  $S$ .

## Some References

This talk is based on material published in [1, 2, 3, 4] and some unpublished material.

Subjects related to this talk been studied over the last seven decades. In a very early paper on computable analysis, Specker [17] proves that

$$\mathcal{S}_D \subset \mathcal{S}_{10E} \subset \mathcal{S}_C$$

where  $\mathcal{S}$  is the class of primitive recursive functions,  $\mathcal{S}_{10E}$  is the set of irrationals that have a primitive recursive decimal expansion,  $\mathcal{S}_D$  is the set of irrationals that have a primitive recursive Dedekind cut and  $\mathcal{S}_C$  is the set of irrationals that have a primitive recursive Cauchy sequence (Specker sequences were introduced in the same paper). In addition to Specker's paper there are works by Mostowski [15], Lehman [14], Ko [11, 12], Labhalla & Lombardi [13], Weihrauch [18], Skordev et al. [16], Georgiev [9] and quite a few more.

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