

# SJT- reducibility and its equivalents

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## Main Def

SJT-reducibility was introduced in Exercise of my 2009 book. It weakens  $\leq_T$ .

### Definition (Main: SJT-reducibility)

- ▶ For an oracle  $B$ , a  $B$ -c.e. trace is a u.c.e. in  $B$  sequence  $\langle T_n \rangle_{n \in \mathbb{N}}$  of finite sets.
- ▶ For a function  $h$ , such a trace is  $h$ -bounded if  $|T_n| \leq h(n)$  for each  $n$ .
- ▶ A set  $A$  is jump-traceable if there is a computably bounded  $\emptyset$ -c.e. trace  $\langle T_n \rangle_{n \in \mathbb{N}}$  such that  $J^A(n)$  is in  $T_n$  if it is defined.

For sets  $A, B$ , we write  $A \leq_{SJR} B$  if for each order function  $h$ , there is a  $B$ -c.e.,  $h$ -bounded trace for  $J^A$ .

This is transitive by an easy argument.

## Proposition

For each  $K$ -trivial set  $B$ , there is a c.e.  $K$ -trivial set  $A$  such that  $A \not\leq_{SJR} B$ .

A **Demuth test** is a sequence  $\langle G_m \rangle_{m \in \mathbb{N}}$  of open subsets of  $\{0, 1\}^{\mathbb{N}}$  such that  $\lambda G_m \leq 2^{-m}$  and there is an  $\omega$ -c.a. function  $p$  such that  $G_m = [W_{p(m)}]^\complement$ .

## Definition (Weak Demuth randomness)

- ▶ A **nested Demuth test** is a Demuth test  $\langle G_m \rangle_{m \in \mathbb{N}}$  such that  $G_m \supseteq G_{m+1}$  for each  $m$ .
- ▶ One says that  $Z$  is **weakly Demuth random** if  $Z \notin \bigcap_m G_m$  for each nested Demuth test  $\langle G_m \rangle_{m \in \mathbb{N}}$ .

A cost function is a computable functions  $c : \mathbb{N} \times \mathbb{N} \rightarrow \mathbb{Q}^+$ .  
 $c(x, s)$  is the cost of changing  $A(x)$  at stage  $s$ .  $\langle A_s \rangle$  satisfies  $c$  if the total of changes is finite.

### Definition

Let  $c$  be a cost function. For sets  $A, B$ , we write  $A \models_B c$  if there is a  $B$ -computable enumeration of  $\langle A_s \rangle$  satisfying  $c$ .

### Definition (Benign cost functions)

A cost function  $\mathbf{c}$  is **benign** if from a rational  $\epsilon > 0$ , we can compute a bound on the length of any sequence  $n_1 < s_1 \leq n_2 < s_2 \leq \dots \leq n_\ell < s_\ell$  such that  $\mathbf{c}(n_i, s_i) \geq \epsilon$  for all  $i \leq \ell$ . For example,  $\mathbf{c}_\Omega$  is benign, with the bound being  $1/\epsilon$ .

**Theorem** (see Logic Blog, 2020)

The following are equivalent for  $K$ -trivial c.e. sets  $A, B$ .

- (a)  $A \leq_{SJR} B$
- (b)  $A \models_B \mathbf{c}$  for every benign cost function  $\mathbf{c}$
- (c)  $A \leq_T B \oplus Y$  for each Martin-Löf-random set  $Y$  that is not weakly Demuth random
- (d)  $A \leq_T B \oplus Y$  for each Martin-Löf-random set  $Y \in \mathcal{C}$ , where  $\mathcal{C}$  is the class of the  $\omega$ -c.a., superlow, or superhigh sets.