

National University of Singapore  
School of Computing

## CS2040S - Data Structures and Algorithms

### Midterm Test

(Semester 1 AY2024/25)

Time Allowed: 70 minutes

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#### INSTRUCTIONS TO CANDIDATES:

1. Do **NOT** open this assessment paper until you are told to do so.
2. This assessment paper contains ONE (1) section with a few short questions.  
It comprises TEN (10) printed pages, including this page.
3. This is an **Open Book Assessment**.  
Only non-programmable calculator is allowed in this assessment.
4. Answer **ALL** questions within the **boxed space** of the answer sheet (page 07-10).  
There are a few starred (\*) boxes: free 1 mark if left blank but 0 for wrong answer (no partial).  
The answer sheet is at page 07-10, just hand in those pages.  
You can use either pen or pencil. Just make sure that you write **legibly!**
5. Important tips: Pace yourself! Do **not** spend too much time on one (hard) question.  
Read all the questions first! Some (subtask) questions might be easier than they appear.
6. You can use **pseudo-code** in your answer but beware of penalty marks for **ambiguous answer**.  
You can use **standard, non-modified** classic algorithm in your answer by just mentioning its name, e.g., run Merge Sort on ArrayList  $A$ , etc.
7. The total marks is 70. All the best :)

## A Short Questions (70 marks)

### A.1 Java Code Comprehension and Analysis (15 marks)

You are given the following Java code that correctly solves one programming contest problem in Kattis:

```
import java.util.*;

public class A1 {
    public static void main(String[] args) throws Exception {
        Scanner sc = new Scanner(System.in);

        // input
        int n = sc.nextInt(), m = sc.nextInt(); sc.nextLine();
        ArrayList<Integer> A = new ArrayList<>();
        for (int i = 0; i < n; ++i)
            A.add(sc.nextInt());

        // the computation
        int ans = 0;
        for (int i = 0; i <= n-m; ++i) {
            int something = 0;
            for (int j = 0; j < m; ++j)
                if (A.get(i+j)%2 == 0)
                    ++something;
            if (something >= 2)
                ++ans;
        }

        // output
        System.out.println(ans); // what is actually computed?
    }
}
```

- (2 marks) If we enter 8 3 as the first line of input and then 1 2 3 4 5 6 7 8 as the second line of input, what will this Java code produce?
- (2 marks) If we enter 9 4 as the first line of input and then 1 2 3 5 7 8 9 11 5 as the second line of input, what will this Java code produce?
- (5 marks) Describe, in high-level, what is being computed in the “the computation” section? Any logical explanation will be accepted.

4. (3 marks) What is the tightest Big O time complexity of this solution, in terms of  $n$  and  $m$ ?

PS: You are given additional information that  $1 \leq m < n$ .

5. (3\* marks) ~~Can you~~ Please rewrite the Java code so that it solves the same problem, but in a faster (i.e., tighter) Big O time complexity (*the time complexity should **not** involve variable  $m$* ).

## A.2 Big-O Time Complexity Analysis (10 marks)

There is an unknown algorithm that process a Java ArrayList  $A$  (containing  $N$  integers). This algorithm has been correctly analyzed to have time complexity of  $O(N \log N)$ . There is no randomized component in this unknown algorithm.

For the five statements below, which statement(s) is/are **always True for all cases**|**can be False on at least one case** given the information above (1 mark per correct decision)? Give a short one sentence explanation for your decision (1 mark for each logical explanation).

1. If  $N = 2^{17} = 131\,072$ , then the total number of operations is **exactly** 2 228 224 operations.
2. This algorithm can be improved to  $O(N)$ .
3. We can also say that this algorithm runs in  $O(N^2)$ .
4. If we change ArrayList  $A$  into PriorityQueue  $PQ$ , the algorithm will run in  $O(N)$ .
5. *Asymptotically, as  $N$  becomes very large*, there is no test-case can make this algorithm executes  $N^2$  operations or more.

## A.3 Another Merge Sort Variant (5 marks)

In class, you learn about the following standard Merge Sort algorithm:

```
method mergeSort(array A, integer low, integer high)
    // the array to be sorted is A[low..high]
    if (low < high) // base case: low >= high (0 or 1 item)
        int mid = (low+high) / 2
        mergeSort(a, low, mid) // divide into two halves
        mergeSort(a, mid+1, high) // then recursively sort them
        merge(a, low, mid, high) // conquer: the O(n) merge subroutine
```

Suppose, you modify it into the following

```
method mergeSort(array A, integer low, integer high)
    // the array to be sorted is A[low..high]
    if (low < high) // base case: low >= high (0 or 1 item)
        int mid = (low+high) / 2
        mergeSort(a, low, mid) // divide into two halves
        mergeSort(a, mid+1, high) // then recursively sort them
        call an O(n log n) sort to sort A[low..high] // ← the change
        // can be anything, perhaps call Arrays.sort(A, low, high)
        // that only sort this subarray A[low..high]
```

- (3 marks) Just choose one of the following (2 marks):  
This modification of Merge Sort is [faster|no different|slower]  
than the original version of Merge Sort discussed in class.  
Give a short one sentence explanation for your decision (1 mark)
- (2\* mark) What is the time complexity of this modification of Merge Sort?

#### A.4 Analyze These Statements (20 marks)

For each of the five statements below, choose whether it is **correct|incorrect** (1 mark). If you say it is correct, write a short one sentence explanation for your decision. If you say it is incorrect, just show one counterexample (the other 1 mark).

- This Java code runs in worst-case time complexity of  $O(n^2)$  and this analysis is tight.

```
int counter = 0;
for (int i = n; i >= 1; --i)
    for (int j = 1; j <= n/i; ++j)
        ++counter;
```

- (Randomized) Quick sort is the *best* sorting algorithm to sort *any* set of  $n$  Integers.
- Java `Collections.sort` can run *slower* than  $O(n \log n)$  (and more than  $O(n^2)$ ) if we use it to sort a collection of  $n$  Java Strings (assume each String has  $m$  Characters).
- We *can* compute the number of inversions (a.k.a. number of Bubble Sort swaps) of an ArrayList  $A$  with  $n$  integers faster than  $O(n^2)$ .
- Singly Linked List discussed in Lecture 4a+4b is a *better* data structure to implement the Stack ADT `push(key)` and `last.key = pop()` than using Java ArrayList, i.e., a resizable array.
- Doubly Linked List discussed in Lecture 4a+4b is a *better* data structure to implement the Deque ADT `offerFirst(key)`, `offerLast(key)`, `key1 = pollFirst()`, and `key2 = pollLast()` than using Java ArrayList, i.e., a resizable array. Important: For this question, the maximum number of elements in the Deque is already known beforehand.
- A Stack/Queue is a Last-In-First-Out/First-In-First-Out data structure, respectively. Thus, there is *no way* we can insert a sequence of  $n$  ( $n > 1$ ) Integers into a Stack and (the same sequence of  $n$  Integers into) a Queue and when we peek the top/front of the Stack/Queue, respectively, we see the same Integer (value).
- Suppose that you have two Linked Lists  $L_a$  and  $L_b$ .  $L_a/L_b$  contains  $n/m$  unsorted alphabets  $['A'..'Z']$ , respectively ( $100 < n, m < 10\,000; n \neq m$ ). As we cannot get an element at index  $i$  in  $O(1)$  if we use Linked List, to check whether an alphabet is inside *both*  $L_a$  and  $L_b$ , we need an  $O(n \times m)$  algorithm that is roughly like this:

```

for (Character a: La)
  ..for (Character b: Lb)
    ...if (a.equals(b)) return true;
return false;

```

9. The smallest element in a Binary *Max* Heap is always located at the bottommost, rightmost leaf, i.e., at index  $n$  of a compact array  $A$  that describes the Binary *Max* Heap.
10. Suppose that we need to use a *special kind* of ADT Priority Queue where all enqueue (Insert( $v$ )) operations will be performed first *before* all subsequent dequeue (ExtractMax()) operations (from highest priority to lowest priority). For this kind of ADT PQ, we *can* use another data structure and/or algorithm to achieve similar time complexities as if we use Binary (Max) Heap.

### A.5 An Algorithm that uses Priority Queue (15 marks)

You are not yet told what this algorithm does, but it is from a real algorithm that you will eventually learn in a future algorithm course (if there is no change in curriculum).

```
import java.util.*;
```

```

public class A5 {
  public static void main(String [] args) {
    int n = 5;
    int [] charFreq = { 5, 1, 6, 10, 3 };

    ArrayList<Integer> ds = new ArrayList<>();
    for (int i = 0; i < n; ++i)
      ds.add(charFreq[i]);

    while (ds.size() > 1) {
      Collections.sort(ds);
      Integer x = ds.get(0);
      ds.remove(0);
      Integer y = ds.get(0);
      ds.remove(0);
      ds.add(x+y);
      System.out.println(x.toString() + "+" + y.toString() + "=" + (x+y));
    }
  }
}

```

1. (4 marks) Output 4 lines that will be the output of that code, if executed.
2. (3 marks) If we change  $n = 7$  and  $charFreq = \{1, 2, 3, 4, 5, 6, 7\}$  What will be the *last line* of the output?

3. (3 marks) What is the time complexity of that code in terms of  $n$ ?

You need to assume that  $n$  can be a very big number, not just  $n = 5$  in this example.

4. (5\* marks) ~~Can you~~ Please rewrite the Java code so that it solves the same problem, but in a faster (i.e., tighter) Big O time complexity. Hint: Do not call `Collections.sort` repeatedly.

There is a better data structure for this problem.

### A.6 The Last Question (5 marks)

To qualify for up to easy 5 marks, you need to **write both full names correctly**.

My CS2040S lecturer is ..... and Teaching Assistant (TA) is .....

Write a **short** (maybe limit yourself to around 2 minutes to do this and about 3-4 sentences) **but honest (and not anonymous)** feedback on what you have experienced in the first 6 weeks of CS2040S in Semester 1 AY 2024/25 (including Week -02/-01 experience, if any). Feedback that are shared by *majority* (**not a one-off**) and can be easily incorporated (e.g., Prof Halim, do not travel again on Week 09 like Week 04 and 06 is very hard to change) to make the next 7 weeks of CS2040S better will be done. Grading scheme: 0-blank, 3-considered trivial feedback but not blank, 5-good and constructive feedback, thanks. (Penalty -1 mark for each wrong name above...).

# The Answer Sheet

Write your Student Number in the box below using **(2B) pencil**.

**Do NOT write your name.**

STUDENT NUMBER									
A									
U	<input type="radio"/>	0	0	0	0	0	0	0	A N
A	<input checked="" type="radio"/>	1	1	1	1	1	1	1	B R
HT	<input type="radio"/>	2	2	2	2	2	2	2	E U
NT	<input type="radio"/>	3	3	3	3	3	3	3	H W
		4	4	4	4	4	4	4	J X
		5	5	5	5	5	5	5	L Y
		6	6	6	6	6	6	6	M
		7	7	7	7	7	7	7	
		8	8	8	8	8	8	8	
		9	9	9	9	9	9	9	<input type="checkbox"/>

Box A.1.1 Check Your Understanding 1 (write one integer)

Box A.1.2 Check Your Understanding 2 (write one integer)

Box A.1.3 Explain that Java code in high-level

Box A.1.4 What is the time complexity in terms of  $n$  and  $m$ ?

Box A.1.5\* (1 if blank, 0 if wrong) Design a faster algorithm, regardless of  $m$

Box A.2 always True|can be False; and why? (5 statements)

  
  
  
  

Box A.3.1. faster|no different|slower; and why?



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Box A.3.2\* (1 if blank, 0 if wrong) What is the time complexity?

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Box A.4 1st-5th boxes, correct|incorrect; and why? (first 5 statements)


Box A.4. 6th-10th boxes, correct|incorrect; and why? (next 5 statements)


Box A.5.1 The output of the given code

Box A.5.2 The output of the next test case

Box A.5.3 The time complexity in terms of  $n$

Box A.5.4\* (1 if blank, 0 if wrong) Design a faster algorithm that does not sort each time

Box A.6 The Last Question: Lecturer and TA name, plus honest feedback

– END OF PAPER; All the Best –