

**CS2100 Assignment #2**  
AY2024/25 Semester 1  
**Deadline: Monday, 14 October 2024, 1:00pm**  
TEMPLATE FOR SUBMISSION (ANSWERS)

Full name:

Tutorial grp: T

**Q1.** (Total: 15 marks)

Cycle time:  ps [4 marks]

Clock frequency:  GHz [3 marks]

Time taken for `beq` instruction:  ps [3 marks]

Optimization: new  [5 marks]

Explain your answers below.

Cycle time is based on the slowest instruction. That is `sw` and `lw` in this case.

`sw` → Fetch (5 ps) + RegFile (2 ps) + ALU (4 ps) + MemWrite (7 ps) = 18 ps

`lw` → Fetch (5 ps) + RegFile (2 ps) + ALU (4 ps) + MemRead (5 ps) + RegFile (2 ps) = 18 ps

Clock frequency =  $(1/18) * 10^{12}$  Hz = **55 GHz**

`beq` → Fetch (5 ps) + RegFile (2 ps) + ALU (4 ps) = **11 ps**

For optimisation, we only look at the slowest instructions.

For `lw`, RegFile reduces by 2 ps and ALU by 2 ps, so max reduction is 2 ps using either.

For `sw`, RegFile reduces by 1 (no write back) and ALU by 2 ps, so max reduction is 2 ps using ALU.

Hence, going with a new ALU is overall better.

**Q2.** (Total: 5 marks)

Easiest would be to take a line from the PC value and then send it to MemToReg, but this would need changes to MemToReg to take a 3<sup>rd</sup> input line. This would mean that we would have the following control signals:

- RegDst = 0
- RegWrite = 1
- ALUSrc = X
- PCSrc = 0
- ALUop = X
- MemRead and MemWrite are both 0
- MemToReg = 2**

**Q3.** (Total: 3 marks)

(a)  $M_{31} =$   $A + B' + C' + D' + E' + F'$  [1 mark]

(b)  $m_{29} \cdot M_{31} =$   $A' \cdot B \cdot C \cdot D \cdot E' \cdot F$  [2 marks]

Note that given a minterm  $m_x$  and a maxterm  $M_y$  on a function with the same  $n$  variables, where  $x, y \in \{0, 1, 2, 3, 4, 5, \dots, 2^n - 1\}$ , if  $x = y$ , then  $m_x \cdot M_y = 0$ , otherwise  $m_x \cdot M_y = m_x$ .

**Q4.** (Total: 4 marks)

Exercise: prove this.

(a)  $F \cdot G' = \sum m( \mathbf{7, 13, 15} )$  [2 marks]

(b)  $G' \oplus H = \sum m( \mathbf{1, 5, 8, 12} )$  [2 marks]

**Q5.** (Total: 3 marks)

Draw your circuit below.

or

$$\begin{aligned}
 R &= (A \text{ XOR } C) \cdot (B \text{ XNOR } D) \\
 &= (A' \cdot C + A \cdot C') \cdot (B' \cdot D' + B \cdot D) \\
 &= A' \cdot C \cdot B' \cdot D' + A' \cdot C \cdot B \cdot D + A \cdot C' \cdot B' \cdot D' + A \cdot C' \cdot B \cdot D \\
 &= A' \cdot B' \cdot C \cdot D' + A' \cdot B \cdot C \cdot D + A \cdot B' \cdot C' \cdot D' + A \cdot B \cdot C' \cdot D \\
 &= m_2 + m_7 + m_8 + m_{13}
 \end{aligned}$$

**Q6.** (Total: 7 marks)

(a) Number of PIs in the K-map of Z: **6** *(B · C, A' · B', B' · C', B' · D, A' · C and C · D.)* [1 mark]

(b) Number of EPIs in the K-map of Z: **1** *(B · C)* [1 mark]

(c) Number of distinct simplified SOP expressions for Z: **3** [1 mark]

(d) One simplified SOP expression for Z: [2 marks]

***B · C + A' · B' + B' · D*** or ***B · C + A' · B' + C · D***  
or ***B · C + B' · D + A' · C***

(e) One simplified POS expression for Z: [2 marks]

***(B' + C) · (A' + B + D)***

**Q7.** (Total:3 marks)

(a) [1 mark]

<i>A</i>	<i>B</i>	<i>C</i>	<i>D</i>	<i>IsZero</i>
0	0	0	0	<b>1</b>
0	0	0	1	<b>0</b>
0	0	1	0	<b>X</b>
0	0	1	1	<b>0</b>
0	1	0	0	<b>0</b>
0	1	0	1	<b>1</b>
0	1	1	0	<b>X</b>
0	1	1	1	<b>0</b>
1	0	0	0	<b>X</b>
1	0	0	1	<b>X</b>
1	0	1	0	<b>X</b>
1	0	1	1	<b>X</b>
1	1	0	0	<b>0</b>
1	1	0	1	<b>0</b>
1	1	1	0	<b>X</b>
1	1	1	1	<b>1</b>

(b) Simplified SOP expression [2 marks]

*IsZero* = ***A · C + B' · D' + A' · B · C' · D***

## Workings

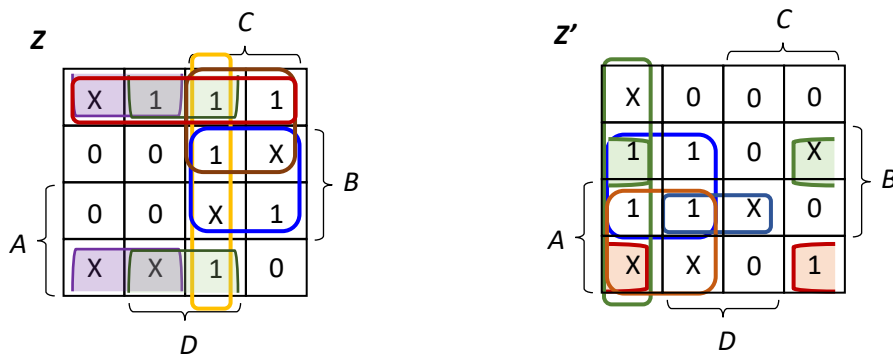
Write your workings here. They will not be graded, but the grader might look at it to figure out where you went wrong.

## Workings for Q3

(b)

$$\begin{aligned}
 m29 \cdot M31 &= M31 \cdot m29 \text{ (commutative law)} \\
 &= (A + B' + C' + D' + E' + F') \cdot (A' \cdot B \cdot C \cdot D \cdot E' \cdot F) \\
 &= A \cdot (A' \cdot B \cdot C \cdot D \cdot E' \cdot F) + B' \cdot (A' \cdot B \cdot C \cdot D \cdot E' \cdot F) \\
 &\quad + C' \cdot (A' \cdot B \cdot C \cdot D \cdot E' \cdot F) + D' \cdot (A' \cdot B \cdot C \cdot D \cdot E' \cdot F) \\
 &\quad + E' \cdot (A' \cdot B \cdot C \cdot D \cdot E' \cdot F) + F' \cdot (A' \cdot B \cdot C \cdot D \cdot E' \cdot F) \text{ (distributive law)} \\
 &= 0 + 0 + 0 + 0 \text{ (complement law } X \cdot X' = 0) \\
 &\quad + E' \cdot (A' \cdot B \cdot C \cdot D \cdot E' \cdot F) + F' \cdot (A' \cdot B \cdot C \cdot D \cdot E' \cdot F) ( \\
 &= (A' \cdot B \cdot C \cdot D \cdot E' \cdot F) + (A' \cdot B \cdot C \cdot D \cdot E' \cdot F) \text{ (idempotent)} \\
 &= (A' \cdot B \cdot C \cdot D \cdot E' \cdot F) \text{ (idempotent)}
 \end{aligned}$$

## K-map for Q6



Working: SOP expression for  $Z'$  is  $B \cdot C' + A \cdot B' \cdot D'$  (see K-map of  $Z'$  above.)

$$Z' = B \cdot C' + A \cdot B' \cdot D'$$

$$Z = (B \cdot C' + A \cdot B' \cdot D')' = (B \cdot C')' \cdot (A \cdot B' \cdot D')' = (B' + C) \cdot (A' + B + D).$$